



Braid ordering: history and connections with knots

Patrick Dehornoy

Laboratoire de Mathématiques Nicolas Oresme
Université de Caen, France

- An introduction to some of the many aspects of the standard **braid order**, with an emphasis on the known connections with **knot theory**.



Conference **ILD**T, Kyoto, May 21, 2015

Plan :

- The Braid Order in Antiquity
- The Braid Order in the Middle Ages
- The Braid Order in Modern Times (Knot Applications)



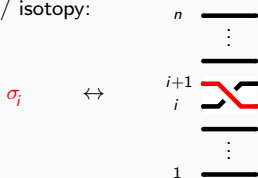
I. The Braid Order in Antiquity: 1985-92

- The set-theoretical roots

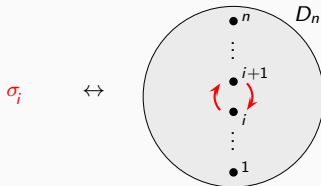
- Definition (Artin 1925/1948): The **braid** group B_n is the group with presentation

$$\langle \sigma_1, \dots, \sigma_{n-1} \mid \begin{array}{l} \sigma_i \sigma_j = \sigma_j \sigma_i \quad \text{for } |i-j| \geq 2 \\ \sigma_i \sigma_j \sigma_i = \sigma_j \sigma_i \sigma_j \quad \text{for } |i-j| = 1 \end{array} \rangle.$$

$\cong \{ \text{braid diagrams} \} / \text{isotopy}:$



\cong mapping class group of D_n (disk with n punctures):



- Braid diagram **colorings**:
 - ▶ start with a set S (“colors”),
 - ▶ apply colors at the left ends of the strands in a braid diagram,
 - ▶ propagate the colors to the right,
 - ▶ compare the initial and final colors.

- Option 1: Colors are preserved in crossings:



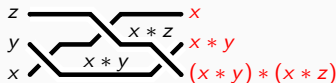
▶ permutation of colors: $B_n \rightarrow \mathfrak{S}_n$

- Option 2: (Joyce, Matveev, Brieskorn 1980s) Colors change under the rule



where $*$ is a (fixed) binary operation on S .

- For an action of B_n on S^n , one needs **compatibility** with the braid relations:



- Fact: One obtains an action of B_n^+ iff $*$ satisfies the **left self-distributivity law (LD)**:

$$x * (y * z) = (x * y) * (x * z).$$

- Classical shelves (or LD-systems) (= sets with an operation obeying the LD-law):
 - ▶ $x * y = y$, leads to $B_n \twoheadrightarrow \mathfrak{S}_n$.
 - ▶ $x * y = xyx^{-1}$, leads to $B_n \rightarrow \text{Aut}(F_n)$ (Artin representation)
 - ▶ $x * y = (1-t)x + ty$, leads to $B_n \rightarrow \text{GL}_n(\mathbb{Z}[t, t^{-1}])$ (Bureau representation)

Note: in these examples, $x * x = x$ always holds.

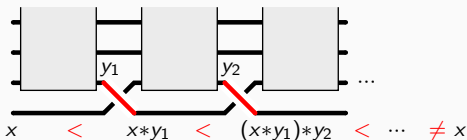
- Definition: A shelf $(S, *)$ is **orderable** if there exists a (left-invariant) linear ordering $<$ on S satisfying $x < x * y$ for all x, y .

Note: (if they exist), orderable shelves are very different: $x < x * x \neq x$.

- Theorem $\frac{1}{2}$: *Orderable shelves exist: free shelves are orderable.*

- Claim: *Theorem 1 (braid order) directly comes from Theorem $\frac{1}{2}$ (orderable LD).*

- Claim: Theorem 1 (braid order) directly comes from Theorem $\frac{1}{2}$ (orderable shelf).
- Ingredient 1 : A σ -positive braid word never represents 1.



- Ingredient 2 : Any two braids are comparable.



Then $\beta <_{\mathcal{D}} \beta'$ iff $(y_1, y_2, \dots) <^{\text{Lex}} (z_1, z_2, \dots)$

- Question: OK, but then, **why** to look for orderable shelves?

- Definition: A **rank** is a set R such that $f : R \rightarrow R$ implies $f \in R$. ??????
- **Assume that** there exists a self-similar set
 (= a set with a nontrivial self-embedding). Then:
 - ▶ There exists a self-similar rank, say R ;
 - ▶ If i, j are self-embeddings of R , then $i : R \rightarrow R$ and $j \in R$, hence
 we can **apply** i to j , obtaining $i(j)$;
 - ▶ “Being a self-embedding” is definable from \in , so $i(j)$ is a self-embedding:
 (“application” is a binary operation on self-embeddings of R).
 - ▶ “Being the image of” is definable from \in , so $\ell=j(k)$ implies $i(\ell)=i(j)(i(k))$,
 that is, $i(j(k)) = i(j)(i(k))$: the “application” operation satisfies the LD law.

- Proposition: If j is a self-embedding of a self-similar rank, then $\text{Iter}(j)$ is a shelf.

\uparrow
 closure of $\{j\}$ under application: $j(j), j(j)(j)\dots$

- Question: Why care about $\text{Iter}(j)$ and prove the previous propositions?

- Theorem 0 (D. 1986): *If j is a self-embedding of a self-similar rank, then the LD-structure of $\text{Iter}(j)$ implies Π_1^1 -determinacy.*

meaning: “ the shelf $\text{Iter}(j)$ is not trivial ”

- Thus: a continuous path from Theorem 0 (about sets) to Theorem 1 (about braids).
- Question: Is the braid order an **application** of Set Theory?
 - ▶ Formally, **no**: braids appear when sets disappear.
 - ▶ In essence, **yes**: orderable shelves have been investigated because Set Theory showed they might exist and be involved in deep phenomena.
- Analogy:
 - ▶ In physics: using **physical** intuition and/or evidence, **guess** some statement, then **pass** it to the mathematician for a formal proof.
 - ▶ Here: using **logical** intuition and/or evidence (\exists self-similar set), **guess** some statement (\exists orderable shelf), then **pass** it to the mathematician for a formal proof.

- The braid order is a complicated object: **non-Archimedean**, **non-Conradian**, ...

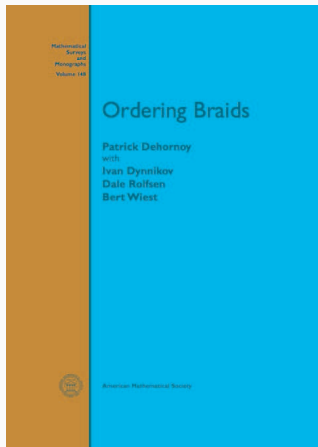
$$\exists \beta, \beta' > 1 \forall p (\beta^p <_D \beta') \quad \exists \beta, \beta' > 1 \forall p (\beta < \beta' \beta^p)$$

- Theorems** (Burckel, D., Dynnikov, Fenn, Fromentin, Funk, Greene, Larue, Rolfsen, Rourke, Short, Wiest, ...):

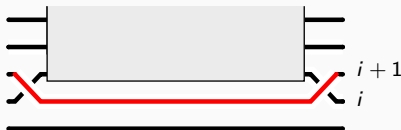
*“Many different approaches
lead to the **same** braid order”.*

- Theorems** (Clay, Dubrovina–Dubrovin, Ito, Navas, Rolfsen, Short, Wiest, ...):

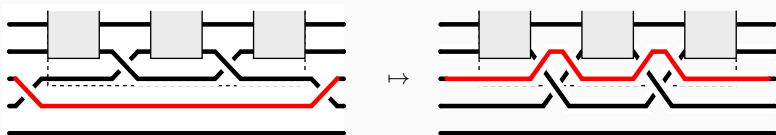
*“There exist many different braid orders
making an interesting **space**”.*



- A σ_i -handle:



- Reducing a handle:



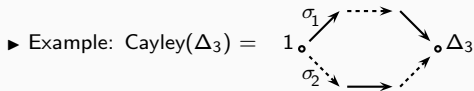
- ▶ Handle reduction is an isotopy;
- ▶ It extends free group reduction;
- ▶ Words with no handle are: the empty word, σ -positive words, σ -negative words.

• Theorem (D. 1995): A braid β satisfies $\beta = 1$ (resp. $\beta > 1$) iff some/any sequence of handle reductions from some/any braid word representing β finishes with **the empty word** (resp. with **a σ -positive word**).

- Aim: Show that there is no infinite sequence of handle reductions.

- **Cayley graph** of B_n : vertices = braids; edge $\beta \xrightarrow{\sigma_i} \beta'$ for $\beta\sigma_i = \beta'$.

- **Cayley(Δ_n^d)**: restriction of the Cayley graph of B_n to the divisors of Δ_n^d
(in the sense of the monoid B_n^+)



- Braid word **drawn in** Cayley(Δ_n^d) from some prescribed vertex:

$\sigma_1\sigma_2\sigma_2^{-1}$ is drawn from 1 in Cayley(Δ_3), but $\sigma_1\sigma_1$ is not.

- Lemma: (i) Every n -strand braid word is drawn in Cayley(Δ_n^d) for $d \gg 0$.
(ii) For every β , the words drawn from β in Cayley(Δ_n^d) are closed under handle reduction.

- Hence: In a sequence of handle reductions,
all words remain drawn in some **finite** fragment of the Cayley graph of B_n .

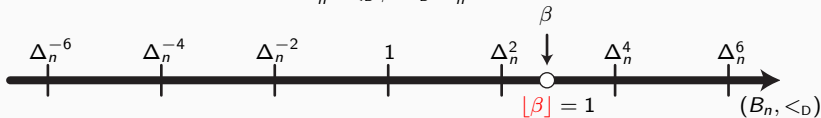


III. The Braid Order in Modern Times: 2000-...

- The floor (after Malyutin–Netstvetaev and Ito)
- Conjugacy via the μ function

- Definition: For β in B_n , the **floor** $\lfloor \beta \rfloor$ is the unique m satisfying

$$\Delta_n^{2m} \leq_D \beta <_D \Delta_n^{2m+2}.$$



- Proposition (Malyutin–Netsvetsev, 2000):

(i) The floor is a quasi-character with defect 1 on B_n : $|\lfloor \beta\gamma \rfloor - \lfloor \beta \rfloor - \lfloor \gamma \rfloor| \leq 1$.

(ii) If β and β' are conjugate, then $|\lfloor \beta \rfloor - \lfloor \beta' \rfloor| \leq 1$.

- Corollary: The **stable floor** $\lfloor \beta \rfloor_s = \lim_p \lfloor \beta^p \rfloor / p$ is the only pseudo-character on B_n that is positive on braids $>_D 1$ and is 1 on Δ_n^2 .

- Principle for using the floor in knot theory:

If $|\lfloor \beta \rfloor|$ is large, then the properties of $\hat{\beta}$ can be read from those of β .

- Lemma: If $||[\beta]|| > 1$, then $\widehat{\beta}$ admits *no destabilisation*.

↑
(assuming $\beta \in B_n$) β is conjugate to no braid $\gamma\sigma_{n-1}^{\pm 1}$ with $\gamma \in B_{n-1}$

- Proof: Assume $\beta \sim \gamma\sigma_{n-1}$ with $\gamma \in B_{n-1}$.

▶ Then $\beta \sim \Delta_n \gamma \sigma_{n-1} \Delta_n^{-1} = \text{sh}(\gamma')\sigma_1$, where $\text{sh} : \sigma_i \mapsto \sigma_{i+1}$ for each i .

▶ Now $1 <_{\text{D}} \text{sh}(\gamma')\sigma_1$, since $\text{sh}(\gamma')\sigma_1$ is σ -positive.

▶ And $\text{sh}(\gamma')\sigma_1 <_{\text{D}} \Delta_n^2$, since $\sigma_1^{-1} \text{sh}(\gamma'^{-1}) \Delta_n^2 = \sigma_1^{-1} \Delta_n^2 \text{sh}(\gamma'^{-1})$ is σ -positive.

▶ Hence, $1 <_{\text{D}} \text{sh}(\gamma')\sigma_1 \leq \Delta_n^2$, that is, $\lfloor \text{sh}(\gamma')\sigma_1 \rfloor = 0$.

▶ Hence, $||[\beta]|| \leq 1$. Idem for $\beta \sim \gamma\sigma_{n-1}^{-1}$...

□

- Proposition (Malyutin–Netsvetov, Ito):

(i) If $||[\beta]|| > 1$, then $\widehat{\beta}$ admits *no exchange move*.

(ii) If $||[\beta]|| > 2$, then $\widehat{\beta}$ admits *no flype*.

- **Theorem** (Malyutin–Netsvetaev, 2004).—

If β satisfies $||[\beta]|| > 1$, then $\widehat{\beta}$ is **prime, non-split, and nontrivial**.

- Proof: For χ a pseudo-character on B_n satisfying $\chi|_{B_{n-1}} = 0$, then $|\chi(\beta)| > \text{defect}(\chi)$ implies that $\widehat{\beta}$ is prime. Apply to $[\]_S$. \square

- **Theorem** (Malyutin–Netsvetaev, 2004).— For every n , there exists $r(n)$ such that for every β in B_n with $||[\beta]|| \geq r(n)$, $\widehat{\beta}$ is represented by a **unique conjugacy class** in B_n .

$$\forall \beta, \beta' \in B_n \quad (\widehat{\beta}' \approx \widehat{\beta} \Rightarrow \beta' \sim \beta)$$

- Proof: For each template move M , there exists r s.t. $||[\beta]|| > r(n)$ implies that $\widehat{\beta}$ is not eligible for M .
By the Birman-Menasco MTWS theory, \exists finitely template moves for each n . \square

- (M.-N., 2000) $r(3) \leq 3$; (Matsuda, 2008) $r(4) \leq 4$; (Ito, 2009) $r(3) = 2$.
conjectured (Ito) $r(n) \leq n - 1$ for each n .

- Theorem (Ito, 2012): For every β in B_n :

$$||\beta|| \leq \frac{4 \cdot \text{genus}(\widehat{\beta})}{n+2} - \frac{2}{n+2} + \frac{3}{2} \leq \text{genus}(\widehat{\beta}) + 1.$$

“The closure of a large braid is a complicated knot”

- Theorem (Ito, 2012): If β satisfies $||\beta|| \geq 2$ and $\widehat{\beta}$ is a knot, then

- ▶ β is periodic iff $\widehat{\beta}$ is a torus knot,
- ▶ β is reducible iff $\widehat{\beta}$ is a satellite knot,
- ▶ β is pseudo-Anosov iff $\widehat{\beta}$ is hyperbolic.

False in general: the trefoil knot is the closure of σ_1^3 (periodic), of $\sigma_1\sigma_2\sigma_3\sigma_1\sigma_2$ (reducible), and of $\sigma_1^3\sigma_2^{-1}$ (pseudo-Anosov).

• Theorem (Ito, 2014): *If H is a nontrivial, non-central normal subgroup of B_n , then, for every γ in B_n , the set $\{\widehat{\beta\gamma} \mid \beta \in H\}$ contains **infinitely many** (hyperbolic) knots.*

• Proof (sketch):

- ▶ The subgroup H is unbounded with respect to $\langle \cdot \rangle_D$: $\forall \gamma \in B_n \exists \beta \in H (\gamma \langle_D \beta)$.
nontrivial: uses the **alternating normal form** of braids...
- ▶ Then $\{\beta\gamma \mid \beta \in H\}$ is also unbounded.
- ▶ Hence, $\{\widehat{\beta\gamma} \mid \beta \in H\}$ contains knots of arbitrarily high genus,
hence certainly infinitely many knots.
- ▶ Moreover, one may assume β pseudo-Anosov, hence $\widehat{\beta}$ hyperbolic. □

• Corollary (Ito, 2014): *Let ρ_1, \dots, ρ_k be non-faithful quantum representations of B_n . Then, for every isotopy type τ , there exist infinitely many hyperbolic knots of type τ on which the invariants derived from ρ_1, \dots, ρ_k agree.*

• Corollary (Ito, 2014): *If the **Burau representation** of B_4 is not faithful, then there exists a nontrivial knot with trivial **Jones polynomial**.*

• Theorem (Laver, 1995): For every braid β and every i , one has $\beta^{-1}\sigma_i\beta >_D 1$.

• Corollary: The restriction of the braid order to B_n^+ is a **well-ordering**.

the submonoid of B_n generated by $\sigma_1, \dots, \sigma_{n-1}$ every nonempty subset has a minimal element

• Definition: For β in B_n^+ , put

$$\mu(\beta) = \min\{\beta' \in B_n^+ \mid \beta' \text{ conjugate to } \beta\}.$$

Useful only if it can be computed...

• Conjecture (D., Fromentin, Gebhardt, 2009): For β in B_3^+ ,

$$\mu(\beta\Delta_3^2) = \sigma_1\sigma_2^2\sigma_1 \cdot \mu(\beta) \cdot \sigma_1^2.$$

...more generally, a reasonable hope of computing μ using the alternating normal form, and its analog for the dual braid monoid (Fromentin's **rotating normal form**).

• If successful for conjugacy, try the same approach for Markov equivalence...

- P. Dehornoy, with I. Dynnikov, D. Rolfsen, B. Wiest, *Ordering braids*
Math. Surveys and Monographs vol. 148, Amer. Math. Soc. (2008)
- A. Malyutin and N. Netsvetaev, Dehornoy's ordering on the braid group and braid moves,
St. Petersburg Math. J. 15 (2004) 437-448.
- T. Ito, Braid ordering and knot genus, *J. Knot Th. Ramif.* 20 (2011) 1311-1323.
- T. Ito, Braid ordering and the geometry of closed braids, *Geom. Topol.* 15 (2011) 473-498.
- T. Ito, Kernel of braid representation and knot polynomial, *Math. Zeitschr.* (2015) to appear
DOI 10.1007/s00209-015-1426-7
- J. Fromentin, Every braid admits a short sigma-definite expression,
J. Europ. Math. Soc. 13 (2011) 1591-1631.

www.math.unicaen.fr/~dehornoy