

## Multifraction reduction and the Word Problem for Artin-Tits groups

Patrick Dehornoy

Laboratoire de Mathématiques Nicolas Oresme Université de Caen

Calais, 22 février 2018

◆□ ▶ ◆□ ▶ ◆ □ ▶ ◆ □ ▶ ◆ □ ▶ ◆ □ ▶

• A new approach to the Word Problem for Artin-Tits groups (and other groups),

- ▶ based on a rewrite system extending free reduction,
- ▶ reminiscent of the Dehn algorithm for hyperbolic groups,
- ▶ proved in particular cases, conjectured in the general case.

# • 1. Reduction of multifractions

- The enveloping group of a monoid
- Free reduction
- A two-step extension: (i) division, (ii) reduction

#### • 2. Artin-Tits monoids I

- The FC case: two theorems
- The general case: three conjectures

#### • 3. Interval monoids (joint with F. Wehrung)

- The interval monoid of a poset
- Examples and counter-examples

#### • 4. Artin-Tits monoids II (joint with D. Holt and S. Rees)

- Padded reduction
- The sufficiently large case

- 1. Reduction of multifractions
  - The enveloping group of a monoid
  - Free reduction
  - A two-step extension: (i) division, (ii) reduction
- 2. Artin-Tits monoids I
  - The FC case: two theorems
  - The general case: three conjectures
- 3. Interval monoids (joint with F. Wehrung)
  - The interval monoid of a poset
  - Examples and counter-examples
- 4. Artin-Tits monoids II (joint with D. Holt and S. Rees)

- Padded reduction
- The sufficiently large case

• For every monoid *M*, there exists a unique group  $\mathcal{U}(M)$  and a unique morphism  $\phi: M \to \mathcal{U}(M)$  s.t. every morphism from *M* to a group factors through  $\phi$ . If  $M = \langle S \mid R \rangle^+$ , then  $\mathcal{U}(M) = \langle S \mid R \rangle$ .

▶ Caution! Even if M is cancellative,  $\phi$  need not be injective: Mal'cev conditions.

• <u>Theorem</u> (Ore, 1933): If M is cancellative and satisfies the 2-Ore condition, then M embeds in  $\mathcal{U}(M)$  and every element of  $\mathcal{U}(M)$  is represented as  $ab^{-1}$  with a, b in M. • "group of (right) fractions for M "

Say that a left-divides b, or b is a right-multiple of a, if ∃x (ax = b). ← a ≤ b
 ▶ 2-Ore condition: any two elements admit a common right-multiple.

• <u>Definition</u>: A gcd-monoid is a cancellative monoid, in which 1 is the only invertible element and any two elements admit a left- and a right-gcd.

a  $\leq$ - and a  $\widehat{\leq}$ -greatest lower bound

• <u>Corollary</u>: If M is a gcd-monoid satisfying the 2-Ore condition, then M embeds in U(M) and every element of U(M) is represented by a unique <u>irreducible</u> fraction.

 $ab^{-1}$  with  $a, b \in M$  and right-gcd(a, b) = 1

then  $\leq$  is a partial order

• <u>Example</u>:  $M = B_n^+$ , the *n*-strand braid monoid; more generally, every Garside monoid.

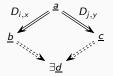
• When the 2-Ore condition fails (no common multiples), no fractional expression.

• Example:  $M = F^+$ , a free monoid; then M embeds in U(M), a free group;

- ▶ No fractional expression for the elements of  $\mathcal{U}(M)$ ,
- But: unique expression a<sub>1</sub>a<sub>2</sub><sup>-1</sup>a<sub>3</sub>a<sub>4</sub><sup>-1</sup> ··· with a<sub>1</sub>, a<sub>2</sub>, ... in M and for i odd: a<sub>i</sub> and a<sub>i+1</sub> do not finish with the same letter, for i even: a<sub>i</sub> and a<sub>i+1</sub> do not begin with the same letter.

a "freely reduced word"

- Proof: (easy) Introduce rewrite rules on finite sequences of positive words:
  - ▶ rule  $D_{i,x} := \begin{cases} \text{for } i \text{ odd, delete } x \text{ at the end of } a_i \text{ and } a_{i+1} \text{ (if possible...)}, \\ \text{for } i \text{ even, delete } x \text{ at the beginning of } a_i \text{ and } a_{i+1} \text{ (if possible...)}. \end{cases}$
  - ▶ Then the system of all rules  $D_{i,x}$  is (locally) confluent:



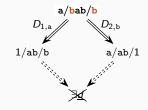
▶ Every sequence <u>a</u> rewrites into a unique irreducible sequence ("convergence").

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

- When M is not free, the rewrite rule  $D_{i,x}$  can still be given a meaning:
  - no first or last letter,
  - ▶ but left- and right-divisors:  $x \leq a$  means "x is a possible beginning of a".

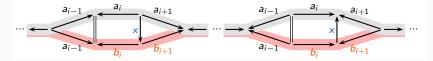
▶ rule  $D_{i,x} := \begin{cases} \text{for } i \text{ odd, right-divide } a_i \text{ and } a_{i+1} \text{ by } x \text{ (if possible...),} \\ \text{for } i \text{ even, left-divide } a_i \text{ and } a_{i+1} \text{ by } x \text{ (if possible...).} \end{cases}$ 

- Example:  $M = B_3^+ = \langle a, b \mid aba = bab \rangle^+$ ;
  - ▶ start with the sequence (a, aba, b), better written a/aba/b ("multifraction")

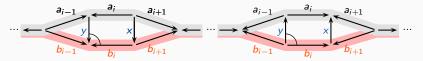


- ▶ no hope of confluence...
- ▶ consider more general rewrite rules.

- Diagrammatic representation of elements of M:  $\xrightarrow{a} \mapsto a$ , and multifractions:  $\xrightarrow{a_1} \xrightarrow{a_2} \xrightarrow{a_3} \dots \mapsto a_1/a_2/a_3/\dots \mapsto \phi(a_1)\phi(a_2)^{-1}\phi(a_3)\dots$  in  $\mathcal{U}(M)$ .
  - ▶ Then: commutative diagram  $\leftrightarrow$  equality in  $\mathcal{U}(M)$ .
- Diagram for  $D_{i,x}$  (division by x at level i): declare  $\underline{a} \cdot D_{i,x} = \underline{b}$  for



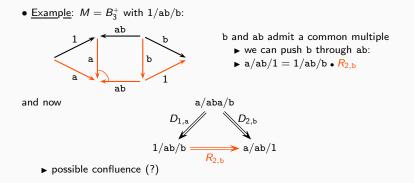
• Relax "x divides  $a_i$ " to "lcm $(x, a_i)$  exists": declare  $\underline{a} \cdot R_{i,x} = \underline{b}$  for



 <u>Definition</u>: "<u>b</u> obtained from <u>a</u> by reducing x at level i": <u>divide</u> a<sub>i+1</sub> by x, push x through a<sub>i</sub> using lcm, <u>multiply</u> a<sub>i-1</sub> by the remainder y.

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

- <u>a</u> D<sub>i,x</sub> is defined if x divides both a<sub>i</sub> and a<sub>i+1</sub>;
  <u>a</u> R<sub>i,x</sub> is defined if x divides a<sub>i+1</sub>, and x and a<sub>i</sub> have a common multiple.
- <u>Fact</u>:  $\underline{b} = \underline{a} \bullet R_{i,x}$  implies that  $\underline{a}$  and  $\underline{b}$  represent the same element in  $\mathcal{U}(M)$ .



• In this way: for every gcd-monoid M, a rewrite system  $\mathcal{R}_M$  ("reduction"), walking among the various multifractions that represent each given element of  $\mathcal{U}(M)$ .

• <u>Theorem</u>: (i) If M is a noetherian gcd-monoid satisfying the 3-Ore condition, then M embeds in  $\mathcal{U}(M)$  and  $\mathcal{R}_M$  is convergent: every element of  $\mathcal{U}(G)$  is represented by a unique  $\mathcal{R}_M$ -irreducible multifraction.

(ii) If, moreover, M is strongly noetherian and has finitely many basic elements, then the word problem for U(M) is decidable.

- ► *M* is noetherian: no infinite descending sequence for left- and right-divisibility.
- ▶ *M* is strongly noetherian: exists a pseudo-length function on *M*. ( $\Rightarrow$  noetherian)
- ► *M* satisfies the 3-Ore condition: three elements that pairwise admit

a common multiple admit a global one. (2-Ore  $\Rightarrow$  3-Ore)

- ▶ right-basic elements: obtained from atoms repeatedly using the right-complement operation:  $(x, y) \mapsto x'$  s.t. yx' = right-lcm(x, y).
- <u>Proof</u>: (i) The rewrite system  $\mathcal{R}_M$  is convergent:
  - ▶ noetherianity of *M* ensures termination;
  - ▶ the 3-Ore condition ensures confluence.
- (ii) Finitely many basic elements provides an upper bound for possible

common multiples, ensuring that  $\Rightarrow$  is decidable.

- 1. Reduction of multifractions
  - The enveloping group of a monoid
  - Free reduction
  - A two-step extension: (i) division, (ii) reduction

#### • 2. Artin-Tits monoids I

- The FC case: two theorems
- The general case: three conjectures
- 3. Interval monoids (joint with F. Wehrung)
  - The interval monoid of a poset
  - Examples and counter-examples
- 4. Artin-Tits monoids II (joint with D. Holt and S. Rees)

- Padded reduction
- The sufficiently large case

- An Artin-Tits monoid: (S | R)<sup>+</sup> such that, for all s, t in S, there is at most one relation s... = t... in R and, if so, the relation has the form stst... = tsts..., both terms of same length (a "braid relation").
- <u>Proposition</u> (Brieskorn–Saito, 1971): An Artin-Tits monoid satisfies the 2-Ore condition iff it of spherical type.

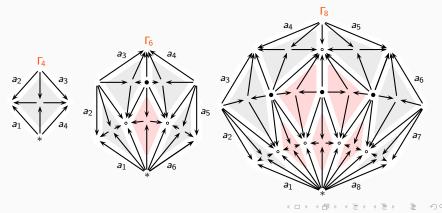
adding  $s^2 = 1$  for every s in S yields a finite Coxeter group

- ▶ "Garside theory"
- <u>Proposition</u>: An Artin-Tits monoid satisfies the 3-Ore condition iff it of FC ("flag complex") type. if  $\forall s, t \in S' \subseteq S \exists s... = t...$  in R, then  $\langle S' \rangle$  is spherical
- <u>Theorem 1</u>: If M is a FC-type Artin-Tits monoid, then every element of the group  $\mathcal{U}(M)$  is represented by a unique  $\mathcal{R}_M$ -irreducible multifraction.
  - ▶ a new normal form

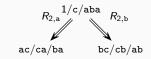
( $\neq$  the Niblo-Reeves n.f. deduced from the action on a CAT(0)-complex)

• <u>Theorem 2</u>: For every n, there exists a universal sequence of integers U(n) such that, whenever M is a FC-type Artin-Tits monoid and <u>a</u> is any depth n multifraction representing 1 in U(M), then <u>a</u> reduces to <u>1</u> by maximal reductions at levels U(n).

- Example: U(8) = (1, 2, 3, 4, 5, 6, 7, 1, 2, 3, 4, 5, 1, 2, 3, 1).
- ▶ Works for every gcd-monoid satisfying the 3-Ore condition.
- ▶ Geometric interpretation: existence of a universal van Kampen diagram  $\Gamma_n$ :

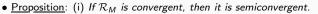


- Every AT-monoid satisfies some of the assumptions:
  - ▶ is strongly noetherian (relations preserve the length of words),
  - ▶ has finitely many basic elements (D.-Dyer-Hohlweg, 2015),
  - ▶ but does not necessarily satisfy the 3-Ore condition, i.e., is of FC-type...
- Example: type A<sub>2</sub>:  $(a, b, c \mid aba = bab, bcb = cbc, cac = aca)^+$ 
  - ▶ the elements a, b, c pairwise admit common multiples, but no global multiple
  - ▶ the rewrite system  $\mathcal{R}_M$  is <u>not</u> confluent:



• <u>Definition</u>:  $\mathcal{R}_M$  is semiconvergent if <u>a</u> represents 1 in  $\mathcal{U}(M)$  iff <u>a</u>  $\Rightarrow^* \underline{1}$ .

▶ Equivalently: conjunction of  $\underline{a} \Rightarrow^* \underline{1}$  and  $\underline{a} \Rightarrow^* \underline{b}$  implies  $\underline{b} \Rightarrow^* \underline{1}$ :



(ii) If M is a strongly noetherian gcd-monoid with finitely many basic elements and  $\mathcal{R}_M$  is semiconvergent, then the word problem for  $\mathcal{U}(M)$  is decidable.

• Conjecture A: For every Artin-Tits monoid M, the system  $\mathcal{R}_M$  is semiconvergent.

- ▶ Would imply the decidability of the word problem for AT groups.
- ▶ Similarity with the Dehn algorithm: no introduction of pairs  $ss^{-1}$  or  $s^{-1}s$ .
- ▶ Supported by massive computer experiments.
- Example: type  $A_2$ :  $(a, b, c \mid aba = bab, bcb = cbc, cac = aca)^+$ 
  - $\blacktriangleright$  We have seen ac/ca/ba  $\iff$  1/c/aba  $\implies$  bc/cb/ab
  - ▶ The quotient ac/ca/ba/ab/cb/bc represents 1 in U(M), hence should reduce to 1:

$$\begin{array}{rcl} \operatorname{ac/ca/ba/ab/cb/bc} & \Rightarrow & \operatorname{ac/cac/b/1/cb/bc} & \operatorname{via} R_{3,ab} \\ & \Rightarrow & \operatorname{ac/cac/bcb/1/1/bc} & \operatorname{via} R_{4,cb} \\ & \Rightarrow & \operatorname{ac/cac/bcb/bc/1/1} & \operatorname{via} R_{5,bc} \\ & \Rightarrow & 1/c/bcb/bc/1/1 & \operatorname{via} R_{1,ac} \\ & \Rightarrow & bc/1/1/bc/1/1 & \operatorname{via} R_{2,cbc} \\ & \Rightarrow & bc/bc/1/1/1/1 & \operatorname{via} R_{3,bc} \\ & \Rightarrow & 1/1/1/1/1/1 & \operatorname{via} R_{1,bc} \end{array}$$

• Finite approximation : *n*-semiconvergence := semiconvergence

restricted to depth n multifractions;

▶  $\mathcal{R}_M$  is 2-semiconvergent iff  $M \subseteq \mathcal{U}(M)$ , which is true (L. Paris, 2001).

• Remember: for FC type, <u>a</u> represents 1 implies <u>a</u> reduces to <u>1</u> by the universal recipe.

• <u>Conjecture</u> B: For every Artin-Tits monoid M, if a multifraction <u>a</u> represents 1 in U(M), then <u>a</u> reduces to <u>1 by the universal recipe</u>.

- ▶ Implies Conjecture A, hence the decidability of the word problem for the group.
- ▶ More precise than A, hence (maybe) more difficult to prove, but easier to test.
- ▶ For depth 4 multifractions, Conjecture A equivalent to Conjecture B.
- ▶ Implies the existence of universal van Kampen diagrams.

▶ call this property "Conjecture B<sup>geom</sup>".

- Question: Does Conjecture  $B^{\text{geom}}$  (hence B) imply that  $\mathcal{U}(M)$  is torsion-free?
- Q<u>uestion</u> (McCammond): Does Conjecture B<sup>geom</sup> (hence B) imply that U(M) satisfies the  $K(\pi, 1)$ -conjecture?

- So far: left reduction: pushing factors to the left
  - ▶ Possible to reverse the definition and push factors to the right
  - ▶ A system  $\widetilde{\mathcal{R}}_M$  ("right reduction"), with entirely symmetric properties.
- <u>Definition</u>: The system  $\mathcal{R}_M$  is cross-confluent if



• <u>Conjecture</u> C: For every Artin-Tits monoid M, the system  $\mathcal{R}_M$  is cross-confluent.

- ▶ Implies Conjecture A, hence the decidability of the word problem for the group.
- ▶ Local version (no \* = one step of reduction) is true.
- ▶ Uniform (stronger) version:  $\forall \underline{a} \exists \nabla \underline{a} \forall \underline{b} (\underline{a} \xrightarrow{\sim} \underline{b} \text{ implies } \underline{b} \longrightarrow \nabla \underline{a}).$
- True for FC type, with  $\nabla \underline{a} := \operatorname{red}(\underline{a})$ .

- 1. Reduction of multifractions
  - The enveloping group of a monoid
  - Free reduction
  - A two-step extension: (i) division, (ii) reduction
- 2. Artin-Tits monoids I
  - The FC case: two theorems
  - The general case: three conjectures
- 3. Interval monoids (joint with F. Wehrung)
  - The interval monoid of a poset
  - Examples and counter-examples
- 4. Artin-Tits monoids II (joint with D. Holt and S. Rees)

- Padded reduction
- The sufficiently large case

ヘロト (日) (日) (日) (日) (日) (日)

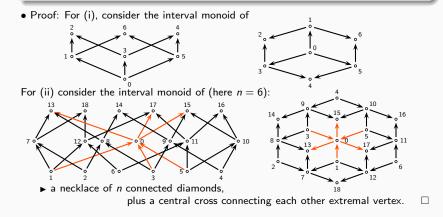
• <u>Definition</u> (F.Wehrung): For  $(P, \leq)$  a poset, the interval monoid of P is  $Int(P) := \langle \{[x, y] \mid x < y \in P\} \mid \{[x, y][y, z] = [x, z] \mid x < y < z \in P\} \rangle^+.$ the intervals of P

• <u>Proposition</u> (D.-Wehrung) If M is the interval monoid of a finite poset P, and M is a gcd-monoid (which is effectively checkable when P is finite), there is a notion of a reducible simple circuit in P such that

- ▶ If all simple circuits of P are reducible, then  $\mathcal{R}_M$  is semiconvergent;
- ▶ If all length  $\leq n$  simple circuits of P are reducible, then  $\mathcal{R}_M$  is n-semiconvergent.

• Checkable conditions for P finite: finitely many simple circuits.

- <u>Proposition</u> (D.-Wehrung): There exist noetherian gcd-monoids such that (i)  $\mathcal{R}_M$  is semiconvergent but not convergent,
  - (ii)  $\mathcal{R}_M$  is n'-semiconvergent for n' < n but not n-semiconvergent.



• The monoids Int(P) are (very) far from Artin-Tits monoids

► A proof of the conjectures must require specific "non-Garside" arguments.

- 1. Reduction of multifractions
  - The enveloping group of a monoid
  - Free reduction
  - A two-step extension: (i) division, (ii) reduction
- 2. Artin-Tits monoids I
  - The FC case: two theorems
  - The general case: three conjectures
- 3. Interval monoids (joint with F. Wehrung)
  - The interval monoid of a poset
  - Examples and counter-examples
- 4. Artin-Tits monoids II (joint with D. Holt and S. Rees)

- Padded reduction
- The sufficiently large case

- Recall:  $\mathcal{R}_M$  semiconvergent if "<u>a</u> represents 1 in  $\mathcal{U}(M)$  implies  $\underline{a} \Rightarrow^* \underline{1}$ ".
- <u>Definition</u>:  $\mathcal{R}_M$  is semiconvergent up to *f*-padding if "<u>a</u> represents 1 in  $\mathcal{U}(M)$  implies  $\underline{1}_{2f(\underline{a})}/\underline{a} \Rightarrow^* \underline{1}$ ". adding  $2f(\underline{a})$  "dummy 1s" at the beginning of <u>a</u>
  - <u>a</u> ⇒\* <u>1</u> implies <u>1</u><sub>2m</sub>/<u>a</u> ⇒\* <u>1</u>.
    but, conversely, adding dummy 1s ("padding") allows for more reductions,
  - ▶ so  $\underline{1}_{2m}/\underline{a} \Rightarrow^* \underline{1}$  need not imply  $\underline{a} \Rightarrow^* \underline{1}$ .

• <u>Proposition</u>: If M is a strongly noetherian gcd-monoid with finitely many basic elements and  $\mathcal{R}_M$  is semiconvergent up to f-padding for some Turing-computable f, then the word problem for  $\mathcal{U}(M)$  is decidable.

• <u>Conjecture</u>  $A^{padded}$ : For every Artin-Tits monoid M, the system  $\mathcal{R}_M$  is semiconvergent up to Turing-computable padding.

▶ Would imply the decidability of the word problem for AT groups.

ヘロト (日) (日) (日) (日) (日) (日)

• An Artin-Tits monoid is said to be of sufficiently large type if, in any triangle in the associated Coxeter diagram,

- either no edge has label 2 ("large type"),
- ▶ or all three edges have label 2 ("right-angled"),
- $\blacktriangleright$  or at least one edge has label  $\infty$  ("free").

• <u>Theorem</u> (D.-Holt-Rees): If M is an AT-monoid of sufficiently large type, then  $\mathcal{R}_M$  is semiconvergent up to a quadratic padding.

• The "first" open case (neither FC nor sufficiently large):

 $\langle a, b, c, d \mid aba = bab, aca = cac, bcb = cbc, ada = dad, bdb = dbd, cd = dc \rangle^+$ .

- <u>P. Dehornoy</u>, Multifraction reduction I: The 3-Ore case and Artin-Tits groups of type FC J. Combinat. Algebra 1 (2017) 185-228, arXiv:1606.08991
- <u>P. Dehornoy</u>, Multifraction reduction II: Conjectures for Artin-Tits groups J. Combinat. Algebra 1 (2017) 229-287, arXiv:1606.08995
- <u>P. Dehornoy</u> & <u>F. Wehrung</u>, *Multifraction reduction III: The case of interval monoids* J. Combinat. Algebra1 (2017) 341-370, arXiv:1606.09018
- <u>P. Dehornoy</u>, <u>D. Holt</u>, <u>& S. Rees</u>, *Multifraction reduction IV: Padding and Artin-Tits groups of sufficiently large type* J. Pure Appl. Algebra, to appear, arXiv:1701.06413
- F. Wehrung, Gcd-monoids arising from homotopy groupoids

hal:01338106

• <u>P. Dehornoy</u>, *MoKa ("monoid calculus")*, MacOS/Docker executable binaries www.math.unicaen.fr/~dehornoy/programs/.